

# Quotient–Comprehension Chains

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Quotients and comprehension are fundamental mathematical constructions that can be described via adjunctions in categorical logic. This paper reveals that quotients and comprehension are related to measurement, not only in quantum logic, but also in probabilistic and classical logic. This relation is presented by a long series of examples, some of them easy, and some also highly non-trivial (esp. for von Neumann algebras). We have not yet identified a unifying theory. Nevertheless, the paper contributes towards such a theory by introducing the new quotient-and-comprehension perspective on measurement instruments, and by describing the examples on which such a theory should be built.

## 1 Introduction

Measurement is a basic operation in quantum theory: the act of observing a quantum system. It is characteristic of the quantum world that such an observation disturbs the system under measurement: it has a side-effect. In [12] a categorical description of measurement is given that takes such side-effects into account. We sketch the essentials, omitting many details. For each predicate  $p$  on a type/object  $A$  in this theory, there is an ‘instrument’ map

$$A \xrightarrow{\text{instr}_p} A + A \quad (1)$$

that performs the act of measuring  $p$ . We write  $A + A$  for the coproduct/sum of  $A$  with itself, which comes equipped with left and right insertion/coprojection maps  $\kappa_1, \kappa_2: A \rightarrow A + A$ . Intuitively, the map  $\text{instr}_p$  gives an outcome in the left summand of  $A + A$  if  $p$  holds, and in the right component otherwise. The side-effect associated with the instrument is the map  $\nabla \circ \text{instr}_p: A \rightarrow A$ , where  $\nabla = [\text{id}, \text{id}]: A + A \rightarrow A$  is the codiagonal. If  $\nabla \circ \text{instr}_p$  is the identity map  $A \rightarrow A$ , one calls  $p$  *side-effect free*. Measurement in a probabilistic setting is side-effect free, but proper quantum measurement is not.

The set-theoretic case may help to understand this instrument map. For each predicate  $p \subseteq A$  one has  $\text{instr}_p(a) = \kappa_1(a)$  if  $a \in P$  and  $\text{instr}_p(a) = \kappa_2(a)$  if  $a \notin P$ . In [12] it is shown that such instrument maps also exist in a probabilistic and in a quantum setting. In the latter case one works in the opposite of the category of  $C^*$ -algebras, with completely positive unital maps. The instrument (1) then has type  $A \times A \rightarrow A$ , and is defined as  $\text{instr}_p(a, b) = \sqrt{p} \cdot a \cdot \sqrt{p} + \sqrt{1-p} \cdot b \cdot \sqrt{1-p}$ . This is the (generalised) Lüders rule, see for instance, in [2, Eq.(1.3)].<sup>1</sup>

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<sup>1</sup>Three notions of measurement (instrument) commonly appear in the literature. *Sharp* or *projective measurement* corresponds to  $\text{instr}_p$  where  $p$  is a projection [17, §2.2.5], and appears in von Neumann’s projection postulate. *POVM measurement* corresponds to arbitrary  $\text{instr}_p$ , although the post-measurement states are usually left out [17, §2.2.6]. *Generalized measurements* capture the different ways the same POVM can be measured [17, §2.2.3]; in the finite dimensional case, every generalized measurement corresponds to a composition  $(\varphi + \psi) \circ \text{instr}_p$ , where  $\varphi$  and  $\psi$  are automorphisms.

The paper [12] lists several requirements for instrument maps (1). The question remained: do these requirements uniquely determine the instrument maps? Put differently: is the presence of these maps a *property* of a category, or *structure*? The current paper does not solve this fundamental problem. But it does uncover the relevance of the logical notions of quotient and comprehension for measurement.

After the formulation of the theory of instruments (1), it became clear (see [4]) that one can also work with *partial* maps  $A \rightarrow A + 1$  and  $A \rightarrow 1 + A$ . The two of them can be combined into a single instrument map  $A \rightarrow A + A$  via a suitable pullback. More importantly, it was noted that in all of the examples the relevant partial map, called ‘assert’ and written as  $\text{asrt}_p: A \rightarrow A$  in the category of partial maps, is a composite of a quotient map  $\xi$  and a comprehension map  $\pi$ , as in:

$$\begin{array}{ccc}
 A & \xrightarrow{\text{asrt}_p} & A \\
 \searrow \xi & & \nearrow \pi \\
 & A/p^\perp \stackrel{(*)}{=} \{A \mid p\} &
 \end{array}
 \tag{2}$$

where  $p^\perp$  is the negation of  $p$ . Such a connection between the fundamental concepts of quotient, comprehension and measurement is fascinating! Quotients and comprehension have a clean description in categorical logic as adjoints (see below for details). Does that lead to instruments as a property? This question remains unsolved, but now takes another form: diagram (2) involves an equality, marked with  $(*)$ , that seems highly un-categorical: adjoints are determined up-to-isomorphism, so having an equality between them is strange. Still this is what we see in all examples, via obvious choices of quotient and comprehension functors. It is not clear if an equality (or isomorphism) between a quotient  $A/p^\perp$  and a comprehension  $\{A \mid p\}$  is property or structure. This is a topic of active research, that requires investigation of many examples. (We have slightly simplified the picture (2) since there is another operation  $[p]$  involved, but that is not essential at this stage; it will be adjusted below.)

This paper is about the following. Once we started looking for quotients and comprehension in the relevant mathematical models we found them everywhere, often in somewhat disguised form. Uncovering familiar constructions, like (co)support for von Neumann algebras, as quotient and comprehension is mathematically relevant on its own. It changes one’s perspective. Thus, the paper only contains examples. Many different examples, each showing that certain constructions are instances of quotient and comprehension. The examples include vector and Hilbert spaces, sets and topological spaces, various Kleisli categories of monads used for probability theory, commutative rings, MV-modules and  $C^*$ -algebras, and finally (non-commutative) von Neumann algebras. The examples point to decomposition of (commutative) mathematical structures as products of quotients and comprehension, like in ring theory, and used for the sheaf theory of commutative rings.

In summary, we think that quotients and comprehension provide a new fruitful perspective on the nature of quantum measurement. This is illustrated here in many examples. We are fully aware that the general, final explanation is lacking at this stage. But such a general theory must be based on a thorough understanding of the examples. That is the focus of the current paper.

This (missing) underlying general theory will bear some resemblance to recent work in (non-Abelian) homological algebra, see in particular [21] (where similar adjunction chains are studied), but also [13, 7]. Part of the motivation is axiomatising the category of (non-Abelian) groups, following [15]. As a result, stronger properties are used than occur in the current setting (for instance the first isomorphism theorem and left adjoints to substitution, corresponding to bifibrations), which excludes not only our motivating example, the category of von Neumann algebras, but also  $\mathcal{Kl}(\mathcal{D})$  and **Sets** to name but two.

## 2 Comprehension and Quotients for Vector Spaces

This section briefly reviews comprehension and quotients for vector spaces. These constructions are fairly familiar. Their categorical description via a chain of adjunctions, as in (3) below, is probably less familiar. This re-description may help to understand similar such chains in the rest of this paper.

We write **Vect** for the category of vector spaces over some fixed field with linear maps between them. Linear subspaces are organised in a category **LSub**. Its objects are pairs  $(V, P)$ , where  $V$  is a vector space and  $P \subseteq V$  is a linear subspace. A morphism  $(P \subseteq V) \rightarrow (Q \subseteq W)$  in **LSub** is a linear map  $f: V \rightarrow W$  that restricts to  $P \rightarrow Q$ , i.e., that satisfies  $P \subseteq f^{-1}(Q)$ . There is then an obvious forgetful functor **LSub**  $\rightarrow$  **Vect**. It is a poset fibration [10], but that does not play a role here. We view **LSub** as a category of linear predicates, over the category **Vect** of linear types.

Interestingly, there is a chain of adjunctions like in (3). The up going functors  $0, 1: \mathbf{Vect} \rightarrow \mathbf{LSub}$  are for falsum and truth respectively. They send a vector space  $V$  to the least  $0(V) = (\{0\} \subseteq V)$  and greatest  $1(V) = (V \subseteq V)$  subspace. There is a comprehension functor  $(V, P) \mapsto P$  that is right adjoint to truth, and a quotient functor  $(V, P) \mapsto V/P$  that is left adjoint to falsum. The outer adjunctions involve (natural) bijective correspondences:

$$\begin{array}{ccc} & \mathbf{LSub} & \\ \text{Quotient} \swarrow & \begin{array}{c} \begin{array}{ccc} \dashv & \dashv & \dashv \\ \curvearrowright & \downarrow & \curvearrowleft \\ 0 & & 1 \end{array} & & \\ \text{Comprehension} \searrow & \mathbf{Vect} & \end{array} & \\ (P \subseteq V) \mapsto V/P & & (P \subseteq V) \mapsto P \end{array} \quad (3)$$

$$\begin{array}{ccc} 1V = (V \subseteq V) & \xrightarrow{f} & (Q \subseteq W) \\ \hline V & \xrightarrow{g} & Q \end{array} \qquad \begin{array}{ccc} (P \subseteq V) & \xrightarrow{f} & (\{0\} \subseteq W) = 0W \\ \hline V/P & \xrightarrow{g} & W \end{array}$$

The second correspondence says that if  $P \subseteq f^{-1}(\{0\}) = \ker(f)$ , then  $f$  corresponds to a map  $V/P \rightarrow W$ . The quotient uses the equivalence relation  $v \sim_P v'$  iff  $v - v' \in P$ .

The category **LSub** is obtained via what is called the ‘Grothendieck construction’. Since we will use it many times in the sequel, we make it explicit. For convenience we restrict it to posets. We write **PoSets** for the category of posets with monotone functions between them.

**Definition 1** Let  $\mathbf{B}$  be a category, with a functor  $F: \mathbf{B} \rightarrow \mathbf{PoSets}^{op}$ . We write  $\int F$  for the category with pairs  $(X, P)$  as objects, where  $X \in \mathbf{B}$  and  $P \in F(X)$ . A morphism  $f: (X, P) \rightarrow (Y, Q)$  is a map  $f: X \rightarrow Y$  in  $\mathbf{B}$  with  $P \leq F(f)(Q)$ . There is an obvious forgetful functor  $\int F \rightarrow \mathbf{B}$ , given by  $(X, P) \mapsto X$  and  $f \mapsto f$ .

The category **LSub** of linear subspaces is obtained via this Grothendieck construction from the functor  $F: \mathbf{Vect} \rightarrow \mathbf{PoSets}^{op}$ , where  $F(V)$  is the poset of linear subspaces of  $V$ , ordered by inclusion; on a linear map  $f: V \rightarrow W$  we get  $F(f): F(W) \rightarrow F(V)$  by inverse image:  $F(f)(Q) = f^{-1}(Q)$ .

The following general observation about the Grothendieck construction is useful.

**Lemma 2** Assume for a functor  $F: \mathbf{B} \rightarrow \mathbf{PoSets}^{op}$ ,

- each ‘fibre’  $F(X)$  has a least element  $0_X$ ;
- each  $F(X)$  also has a greatest element  $1_X$ , and each  $F(f): F(Y) \rightarrow F(X)$  satisfies  $F(f)(1_Y) = 1_X$ .

Then there are functors  $0, 1: \mathbf{B} \rightarrow \int F$ , namely  $0(X) = (X, 0_X)$  and  $1(X) = (X, 1_X)$ , which are left and right adjoints to the forgetful functor  $\int F \rightarrow \mathbf{B}$ .  $\blacksquare$

We briefly sketch the situation for Hilbert spaces, where quotients are given by (ortho)complements. So let **Hilb**  $\leftrightarrow$  **Vect** be the category of Hilbert spaces, with bounded linear maps between them. Mapping a Hilbert space  $V$  to the poset of closed linear subspaces yields a functor **Hilb**  $\rightarrow$  **PoSets**<sup>op</sup>. We write **CLSub** for the resulting Grothendieck completion, with forgetful functor **CLSub**  $\rightarrow$  **Hilb**. Since both  $\{0\} \subseteq V$  and  $V \subseteq V$  are closed, this functor has both a left and right adjoint, by Lemma 2.



The counit of the adjunction between  $\mathbf{1}$  and comprehension for an object  $(Y, Q)$  in  $\int \square$  is the inclusion  $\pi_Q: Q \rightarrow Y$ . It has the following universal property. For every  $f: X \rightarrow Y + 1$  with  $f(X) \subseteq Q \cup \{*\}$  there is a unique map  $\bar{f}: X \rightarrow Q + 1$  such that  $f = \pi_Q \circ \bar{f}$ . The map  $\bar{f}$  will simply be the restriction of  $f$  to a partial map from  $X$  to  $Q$ .

In this situation consider the following (composite) maps, the first two in  $\mathbf{Sets}_{+1}$ , the last one in  $\mathbf{Sets}$ .

$$\begin{array}{ccc}
 P \xrightarrow{\pi_P} X \xrightarrow{\xi_{-P}} P & X \xrightarrow{\xi_{-P}} P \xrightarrow{\pi_P} X & x \xrightarrow{\text{instr}_P} X + X \\
 x \mapsto x \mapsto x & x \mapsto \begin{cases} x & \text{if } x \in P \\ * & \text{if } x \notin P \end{cases} \mapsto \begin{cases} x & \text{if } x \in P \\ * & \text{if } x \notin P \end{cases} & x \mapsto \begin{cases} \kappa_1 x & \text{if } x \in P \\ \kappa_2 x & \text{if } x \notin P \end{cases} \quad (6)
 \end{array}$$

The first map is the identity; the second one is the ‘assert’ map  $\text{asrt}_P$  from the introduction; and the third one is obtained by combining  $\text{asrt}_P$  and  $\text{asrt}_{-P}$  via a suitable pullback. It is the instrument map for measurement, associated with the predicate  $P \subseteq X$ .

$$\begin{array}{ccc}
 \text{Quotient} & \int \square & \text{Comprehension} \\
 (P \subseteq X) \rightarrow \neg P & \begin{array}{c} \curvearrowright \quad \curvearrowleft \\ \neg \quad 0 \quad \neg \\ \downarrow \quad \downarrow \\ \neg \quad 1 \quad \neg \\ \curvearrowleft \quad \curvearrowright \end{array} & (P \subseteq X) \rightarrow P \\
 & \mathbf{Top}_{+1} & 
 \end{array} \quad (7)$$

There are some relatively straightforward variations of the chain of adjunctions in (5). If one replaces the poset  $\mathcal{P}(X)$  of subsets of a set  $X$  by the poset  $\text{Clopen}(X)$  of clopens of a topological space  $X$  one gets a functor  $\square: \mathbf{Top} \rightarrow \mathbf{PoSets}^{\text{op}}$ .

For a continuous function  $f: X \rightarrow Y + 1$  (which corresponds to a continuous partial function  $f: X \rightarrow Y$  with clopen domain) and clopen  $Q \subseteq Y$  we define  $\square(f)(Q) = f^{-1}(Q \cup \{*\})$  as before. (Note that  $f^{-1}(Q \cup \{*\})$  is clopen.) Again one gets a quotient–comprehension chain (7). For a clopen  $P \subseteq X$  the quotient  $\neg P$  and comprehension  $P$  are the same as in the case of sets (5) but now come with a natural topology induced by  $X$ . To see that this works one checks that all maps involved are continuous.

One obtains a similar chain for the category  $\mathbf{Meas}$  of measurable spaces and measurable maps if one replaces the poset  $\mathcal{P}(X)$  of subsets of a set  $X$  by the poset  $\text{Meas}(X)$  of measurable subsets of a measurable space  $X$ .

Let us think some more about the chain for topological spaces. Since a closed subset of a compact Hausdorff space is again compact, we may restrict the chain (7) to the category  $\mathbf{CH}$  of compact Hausdorff spaces and the continuous maps between them. Since  $\mathbf{CH}$  is dual to a whole slew of ‘algebraic’ categories (as opposed to ‘spacial’ such as  $\mathbf{Top}$ ) we get quotient–comprehension chains for (the opposite of) all those categories as well. For example, we get a quotient–comprehension chain for the opposite category of commutative unital  $C^*$ -algebras with unital  $*$ -homomorphisms via Gelfand’s duality (see e.g. [6]), and for the opposite category of unital Archimedean Riesz spaces with Riesz homomorphisms via Yosida’s duality [23]. Interestingly, there are quotient–comprehension chains for ‘algebraic’ categories which do not seem to have a ‘spacial’ counterpart such as the category of commutative rings and homomorphisms, such as the category  $\mathbf{CRng}^{\text{op}}$  of commutative rings and homomorphisms, as we will see in Section 5.

The categories  $\mathbf{Sets}$ ,  $\mathbf{Top}$ ,  $\mathbf{Meas}$ ,  $\mathbf{CH}$  and  $\mathbf{CRng}^{\text{op}}$  are all *extensive* [3]. In fact, any extensive category  $\mathcal{E}$  with final object has a quotient–adjunction chain of which (5) and (7) are instances. In particular, any topos will have a quotient–adjunction chain. In this general setting, the poset of subsets of a set  $X$  is replaced by the poset of *complemented* subobjects of an object  $X$  of  $\mathcal{E}$ . Details will appear elsewhere.

For our next example we write  $\mathcal{P}_*$  for the *nonempty* powerset monad on  $\mathbf{Sets}$ ,  $\mathcal{Kl}(\mathcal{P}_*)$  for its Kleisli category, and  $\mathcal{Kl}(\mathcal{P}_*)_{+1}$  for the Kleisli category of the lift monad on  $\mathcal{Kl}(\mathcal{P}_*)$ . Thus, maps  $X \rightarrow Y$  in  $\mathcal{Kl}(\mathcal{P}_*)_{+1}$  are functions  $X \rightarrow \mathcal{P}_*(Y + 1)$ . They capture non-deterministic computation, with multiple successor states and possibly also non-termination.



the set of fuzzy predicates  $[0, 1]^X$  forms a what is called an effect module, that is, an effect algebra with a  $[0, 1]$ -action (see [12] for details).

A predicate  $p \in [0, 1]^X$  is called *sharp* if  $p^2 = p$ . This means that  $p(x) \in \{0, 1\}$ , so that  $p$  is a Boolean predicate in  $\{0, 1\}^X$ . Equivalently,  $p$  is sharp if  $p \wedge p^\perp = 0$ . For each predicate  $p \in [0, 1]^X$  there is a least sharp predicate  $\lceil p \rceil$  with  $p \leq \lceil p \rceil$ , and a greatest sharp predicate  $\lfloor p \rfloor \leq p$ , namely:

$$\lceil p \rceil(x) = \begin{cases} 0 & \text{if } p(x) = 0 \\ 1 & \text{otherwise.} \end{cases} \quad \lfloor p \rfloor(x) = \begin{cases} 1 & \text{if } p(x) = 1 \\ 0 & \text{otherwise.} \end{cases}$$

It is easy to see that these least sharp and greatest sharp predicates are each others De Morgan duals, that is,  $\lfloor p^\perp \rfloor = \lceil p \rceil^\perp$ . If  $p$  is a sharp, then  $\lfloor p \rfloor = p = \lceil p \rceil$ .

For a function  $f: X \rightarrow \mathcal{D}(Y+1)$  we define  $\square(f): [0, 1]^Y \rightarrow [0, 1]^X$  as:

$$\square(f)(q)(x) = \sum_{y \in Y} f(x)(y) \cdot q(y) + f(x)(*).$$

Since  $f(x) \in \mathcal{D}(Y+1)$  is a distribution, we have  $\sum_{y \in Y} f(x)(y) + f(x)(*) = 1$ , so that  $\square(f)(1) = 1$ . Hence Lemma 2 applies, so that we have a functor  $\int \square \rightarrow \mathcal{Kl}(\mathcal{D})_{+1}$  with falsum 0 as left adjoint, and truth 1 as right adjoint. Recall that a map  $(X, p) \rightarrow (Y, q)$  in  $\int \square$  is a function  $f: X \rightarrow \mathcal{D}(Y+1)$  with  $p(x) \leq \square(f)(q)(x)$  for all  $x \in X$ .

**Proposition 6** *The distribution monad  $\mathcal{D}$  on Sets, used to model probabilistic computation, gives rise to the chain of adjunctions (10) to the right where  $\{X|p\} = \{x \in X \mid p(x) = 1\}$ , and  $X/p = \{X \mid \lceil p^\perp \rceil\} = \{x \mid p(x) \neq 1\}$ .*

$$\begin{array}{ccc} & \int \square & \\ \text{Quotient} & \begin{array}{c} \begin{array}{ccc} \begin{array}{c} \downarrow \lceil \cdot \rceil \\ \text{0} \end{array} & \begin{array}{c} \lceil \cdot \rceil \\ \downarrow \lceil \cdot \rceil \\ \text{1} \end{array} & \begin{array}{c} \lceil \cdot \rceil \\ \downarrow \lceil \cdot \rceil \\ \text{1} \end{array} \\ \text{0} & \begin{array}{c} \lceil \cdot \rceil \\ \downarrow \lceil \cdot \rceil \\ \text{1} \end{array} & \text{1} \\ \text{1} & \begin{array}{c} \lceil \cdot \rceil \\ \downarrow \lceil \cdot \rceil \\ \text{1} \end{array} & \text{1} \end{array} & \text{Comprehension} \\ (p \in [0, 1]^X) \rightarrow X/p & \mathcal{Kl}(\mathcal{D})_{+1} & (p \in [0, 1]^X) \rightarrow \{X|p\} \end{array} \quad (10)$$

**Proof** For a map  $f: 1Y \rightarrow (X, p)$  in  $\int \square$  we have  $f: Y \rightarrow \mathcal{D}(X+1)$  satisfying  $1 \leq \square(f)(p)$ . This means  $1 = (\sum_x f(y)(x) \cdot p(x)) + f(y)(*)$ , for each  $y \in Y$ . Since  $\sum_x f(y)(x) + f(y)(*) = 1$ , this can only happen if  $f(y)(x) \neq 0 \Rightarrow p(x) = 1$ . But then we can factor  $f$  as  $\bar{f}: Y \rightarrow \{X|p\}$  in  $\mathcal{Kl}(\mathcal{D})_{+1}$ , where  $\bar{f}(y) = \sum_{x, f(y)(x) \neq 0} f(y)(x) |x\rangle + f(y)(*) |*\rangle$ .

In the other direction, given a function  $g: Y \rightarrow \mathcal{D}(\{X|p\} + 1)$  we define the map  $\bar{g}: Y \rightarrow \mathcal{D}(X+1)$  as  $\bar{g}(y) = \sum_{x, p(x)=1} g(y)(x) |x\rangle + g(y)(*) |*\rangle$ . Then, for each  $y \in Y$ ,

$$\square(\bar{g})(p)(y) = \sum_{x, p(x)=1} \bar{g}(y)(x) \cdot p(x) + \bar{g}(y)(*) = \sum_{x, p(x)=1} g(y)(x) + g(y)(*) = 1.$$

The quotient adjunction involves the correspondence (11), which works as follows. Given  $f: (X, p) \rightarrow 0Y$  in  $\int \square$ , then  $f: X \rightarrow \mathcal{D}(Y+1)$  satisfies  $p \leq \square(f)(0)$ . This means that  $p(x) \leq$

$\sum_y f(x)(y) \cdot 0(y) + f(x)(*) = f(x)(*)$ , for each  $x \in X$ . We then define  $\bar{f}: X/p \rightarrow \mathcal{D}(Y+1)$  as  $\bar{f}(x) = \sum_y \frac{f(x)(y)}{p^\perp(x)} |y\rangle + \frac{f(x)(*) - p(x)}{p^\perp(x)} |*\rangle$ . This is well-defined, since  $p^\perp(x) \neq 0$  for  $x \in X/p$ .

In the other direction, given  $g: X/p \rightarrow \mathcal{D}(Y+1)$  we define  $\bar{g}: X \rightarrow \mathcal{D}(Y+1)$  as:

$$\bar{g}(x) = \sum_y p^\perp(x) \cdot g(x)(y) |y\rangle + (p(x) + p^\perp(x) \cdot g(x)(*)) |*\rangle.$$

Notice that this extension of  $g$  outside the subset  $\{X \mid \lceil p^\perp \rceil\} \hookrightarrow X$  is well-defined, since if  $x \notin \{X \mid \lceil p^\perp \rceil\}$ , then  $p(x) = 1$ , so  $p^\perp(x) = 0$ , which justifies writing  $p^\perp(x) \cdot g(x)(y)$ . In that case, when  $p(x) = 1$ , we get  $\bar{g}(x) = 1 |*\rangle$ . This  $\bar{g}$  is a morphism  $(X, p) \rightarrow 0Y$  in  $\int \square$ , since  $p \leq \square(\bar{g})(0)$ , that is  $p(x) \leq \bar{g}(x)(*)$ . This follows since  $p^\perp(x) \geq 0$  and  $g(x)(*) \geq 0$  in  $\bar{g}(x)(*) = p(x) + p^\perp(x) \cdot g(x)(*) \geq p(x)$ . ■

$$\begin{array}{ccc} (X, p) & \xrightarrow{f} & 0Y \\ \hline X/p & \xrightarrow{g} & Y \end{array} \quad (11)$$









The quotient functor sends an effect  $p$  of a von Neumann algebra  $\mathcal{A}$  to the set of elements of  $\mathcal{A}$  of the form  $[p^\perp]a[p^\perp]$  which is denoted by  $[p^\perp]\mathcal{A}[p^\perp]$ . We should note that  $[p^\perp]\mathcal{A}[p^\perp]$  is a linear subspace of  $\mathcal{A}$  which is closed under multiplication, involution  $(-)^*$  and is closed in the weak operator topology, so that  $[p^\perp]\mathcal{A}[p^\perp]$  is itself (isomorphic to) a von Neumann algebra. The unit of  $[p^\perp]\mathcal{A}[p^\perp]$  is  $[p^\perp]$  which might be different from the unit of  $\mathcal{A}$ . The unit of the adjunction between quotient and 0 on the effect  $p \in \mathcal{A}$  is the cPNsU-map  $\xi_p: [p^\perp]\mathcal{A}[p^\perp] \rightarrow \mathcal{A}$  which sends  $a$  to  $\sqrt{p^\perp}a\sqrt{p^\perp}$ .

As in the probabilistic example, we can form the following composites in  $\mathbf{W}_{+1}^*$ .

$$\begin{array}{ccc} [p]\mathcal{A}[p] \xleftarrow{\pi_{[p]}} \mathcal{A} \xleftarrow{\xi_{p^\perp}} [p]\mathcal{A}[p] & \mathcal{A} \xleftarrow{\xi_{p^\perp}} [p]\mathcal{A}[p] \xleftarrow{\pi_{[p]}} \mathcal{A} & (15) \\ \sqrt{p}a\sqrt{p} \longleftarrow \sqrt{p}a\sqrt{p} \longleftarrow a & \sqrt{p}a\sqrt{p} \longleftarrow [p]a[p] \longleftarrow a & \end{array}$$

The map on the left is the identity if the predicate  $p$  is sharp. The map on the right is the ‘assert’ map  $\text{asrt}_p$ , which yields, together with  $\text{asrt}_{p^\perp}$  the instrument  $\text{instr}_p: \mathcal{A} \times \mathcal{A} \rightarrow \mathcal{A}$  in  $\mathbf{W}_{+1}^*$  given by

$$\text{instr}_p(a, b) = \sqrt{p}a\sqrt{p} + \sqrt{1-p}b\sqrt{1-p}$$

precisely as in [12]. Hence we see how the instrument map for measurement for von Neumann algebras is obtained via the logical constructions of quotient and comprehension.

**Proofsketch of Proposition 8 (Comprehension)** We must show that given a von Neumann algebra  $\mathcal{A}$ , an effect  $p \in \mathcal{A}$ , and a map  $f: \mathcal{A} \rightarrow \mathcal{B}$  in  $\mathbf{W}_{+1}^*$  with  $f(p) = f(1)$  there is a unique map  $g: [p]\mathcal{A}[p] \rightarrow \mathcal{B}$  in  $\mathbf{W}_{+1}^*$  with  $g([p]b[p]) = f(b)$ . Put  $g(b) = f(b)$ ; the difficulty is to show that  $f([p]b[p]) = f(b)$ . By a variant of Cauchy–Schwarz inequality for the completely positive map  $f$  (see [19], exercise 3.4)

$$\|f(c^*d)\|^2 \leq \|f(c^*c)\| \cdot \|f(d^*d)\| \quad (c, d \in \mathcal{A})$$

we can reduce this problem to proving that  $f([p]) = f(1)$ , that is,  $f([p^\perp]) = 0$ . Since  $[p^\perp]$  is the supremum of  $p^\perp \leq (p^\perp)^{1/2} \leq (p^\perp)^{1/4} \leq \dots$  and  $f$  is normal,  $f([p^\perp])$  is the supremum of the operators  $f(p^\perp) \leq f((p^\perp)^{1/2}) \leq f((p^\perp)^{1/4}) \leq \dots$ , which all turn out to be zero by Cauchy–Schwarz since  $f(p) = f(1)$ . Thus  $f([p^\perp]) = 0$ , and we are done. Again, for more details, see [22].

(*Quotient*) We must show that given a von Neumann algebra  $\mathcal{A}$ , an effect  $p \in \mathcal{A}$ , and a map  $f: \mathcal{B} \rightarrow \mathcal{A}$  in  $\mathbf{W}_{+1}^*$  with  $f(1) \leq p^\perp$ , there is a unique  $g: \mathcal{B} \rightarrow [p^\perp]\mathcal{A}[p^\perp]$  in  $\mathbf{W}_{+1}^*$  such that  $\sqrt{p^\perp}g(b)\sqrt{p^\perp} = f(b)$ . If  $\sqrt{p^\perp}$  is invertible, then we may define  $g(b) = (\sqrt{p^\perp})^{-1}f(b)(\sqrt{p^\perp})^{-1}$ , and this works. The proof is also straightforward if  $\sqrt{p^\perp}$  is pseudoinvertible (=has norm-closed range). The trouble is that in general  $\sqrt{p^\perp}$  is not (pseudo)invertible. However, using the spectral theorem [9] we can find a sequence  $s_n$  (which converges ultraweakly to the (pseudo)inverse if it exists and) for which  $g(b) = \text{uwlim}_n s_n f(b) s_n$  exists and satisfies the requirements. For further details, see [22]. ■

## 7 Conclusions

This paper uncovers a fundamental chain of adjunctions for quotient and comprehension in many example categories of mathematical structures, in particular von Neumann algebras. This in itself is a discovery. Truly fascinating to us is the role that these adjunctions play in the description of measurement instruments in these examples. To our regret we are unable at this stage to offer a unifying categorical formalisation, since in each of the examples there is an equality connecting adjoints which are determined only up-to-isomorphism. To be continued!

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